## FAGNANO'S PROBLEM Addendum

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The problem treated in the accompanying article is this: Given an arbitrary acute-angled triangle PQR, inscribe within it a triangle ABC, with A on side RP, B on side PQ, and C on side QR, having the smallest possible perimeter. The author establishes, using geometrical arguments, that in the optimal configuration, the following triangle similarities must hold (see Figure 1):

 $\triangle ARC \sim \triangle QBC \sim \triangle ABP \sim \triangle QRP$ ,

and then shows, using trigonometry, that these conditions force A, B, C to be the feet of the altitudes of the triangle.

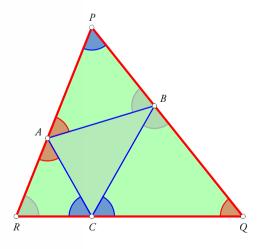
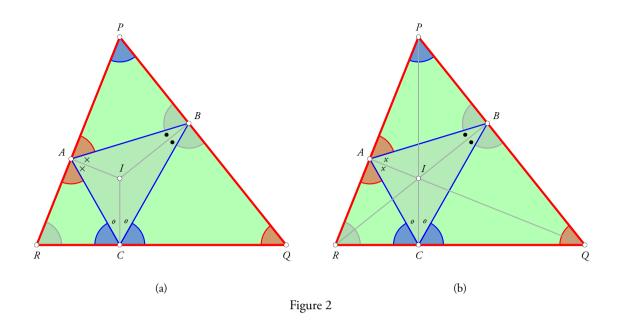


Figure 1. Fagnano's Problem

Here we provide a geometrical proof of this proposition. We also justify the need to impose the condition that triangle PQR should be acute-angled.

*Keywords:* triangle, acute, obtuse, perpendicular, angle bisector, incentre, excentre, collinear



## **Proof of proposition**

Construct the internal bisectors of the angles of  $\triangle ABC$ . The three lines thus constructed meet at the incentre *I* of  $\triangle ABC$ ; see Figure 2 (a).

It is easy to check, using elementary angle computations, that the sides of  $\triangle PQR$  are respectively perpendicular to the three angle bisectors; that is, side QR is perpendicular to the angle bisector CI of  $\measuredangle ACB$ , and so on. But this implies that the sides of  $\triangle PQR$  are respectively the external bisectors of the angles of  $\triangle ABC$  (i.e., side QR is the external angle bisector of  $\measuredangle ACB$ , and so on). This in turn implies that P, Q, R are the ex-centres of  $\triangle ABC$ . And this in turn implies that P, Q, R lie on the (internal) angle bisectors of  $\measuredangle ACB$ ,  $\measuredangle CAB$ ,  $\measuredangle ABC$  respectively. That is, points P, I, C are collinear, as are points Q, I, Aand points R, I, B; see Figure 2 (b).

It follows that *PC*, *QA* and *RB* are the altitudes of  $\triangle PQR$ . This is just what we had set out to prove.  $\Box$ 

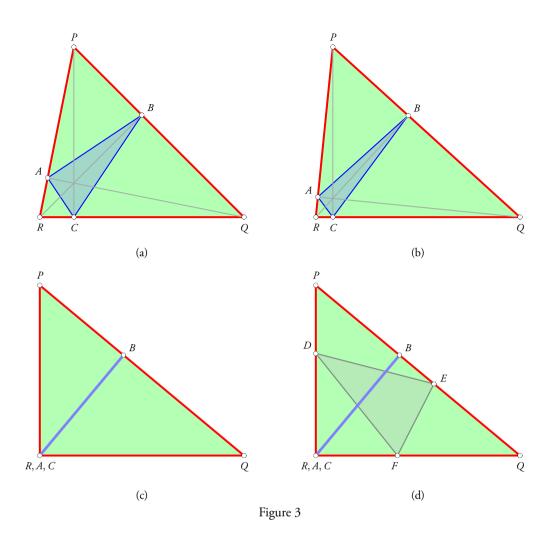
## Why should the triangle be acute angled?

We now justify the need to impose the condition that  $\triangle PQR$  should be acute angled. We accomplish this by considering what happens if  $\triangle PQR$  is right-angled or obtuse-angled.

Figures 3 (a), 3 (b) and 3 (c) show triangles in each of which the angle at vertex R is successively larger than in Figures 1 and 2; it is getting 'closer' to a right angle, and in the limit, Figure 3 (c), the triangle becomes right-angled at vertex R.

Observe carefully what happens: as  $\measuredangle R$  increases, vertices A and C get steadily closer to each other, and in the limit, when the triangle becomes right-angled at vertex R, the two vertices coincide with R. When this happens,  $\triangle ABC$  collapses into segment RB. The configuration will now be as depicted in Figure 3 (c). We infer from this that if  $\triangle PQR$  is right-angled, then the inscribed triangle with least perimeter is a line segment. (Note that in the limiting situation, segment BR is traced out *twice*, which means that the perimeter of  $\triangle ABC$  is twice the length of segment BR.)

It is possible to show directly that if  $\triangle PQR$  is right-angled at *R* and  $\triangle DEF$  is inscribed in  $\triangle PQR$ , then its perimeter cannot be less than twice the length of altitude *RB*. Let *DEF* be any inscribed triangle, as in Figure 3 (d). We now perform the following geometrical operations on this figure: we reflect the



configuration in line *QR* and again in line *PR*. The effect is shown in Figure 4; the two mappings take *E* to points  $E_1$  and  $E_2$  respectively.

We now note the following:

- $E_1F = EF$  and  $E_2D = ED$ , by the very nature of the reflection operation; hence the perimeter of  $\triangle DEF$  is equal to the length of the path  $E_1FDE_2$ .
- $\angle ERE_1 = 2 \angle ERQ$  and  $\angle ERE_2 = 2 \angle ERP$ , so  $\angle E_1RE_2 = 2 \angle PRQ = 180^\circ$ . That is, points  $E_1, R, E_2$  lie in a straight line.
- The length of path  $E_1FDE_2$  is greater than or equal to the length of segment  $E_1E_2$ , i.e., greater than or equal to  $2 \times$  the length of segment *RE*. (This follows from several usages of the result that any two sides of triangle are together greater that the third side.) Hence: perimeter of  $\triangle DEF \ge 2 \times$  the length of segment *RE*.
- The length of segment *RE* is greater than or equal to the length of segment *RB* (because *RB* is perpendicular to *PQ*).
- Hence: perimeter of  $\triangle DEF \ge 2 \times$  the length of segment *RB*.

The stated claim therefore follows: the optimal inscribed triangle ('optimal' in the sense of having the least possible perimeter) is the degenerate triangle consisting of the segment RB traced twice over.

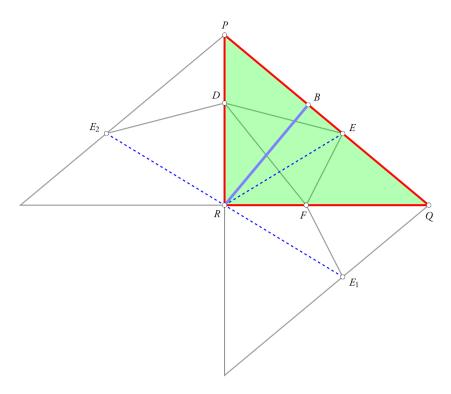


Figure 4

If we continue to increase the size of  $\measuredangle PRQ$ , then we get a triangle which is obtuse-angled at vertex *R*. What is the optimal inscribed triangle in this case? It turns out that we cannot do better than opting for the degenerate triangle which consists of the segment *RB* traced twice over (here, *B* is the foot of the perpendicular from vertex *R* to side *PQ*). We leave the full justification of the statement to you. (Hint: The reflection idea used above will work here as well.)



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